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# Integer sets with prescribed pairwise differences being distinct

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## Abstract

We label the vertices of a given graph  $G$  with positive integers so that the pairwise differences over its edges are all distinct. Let  $\mathcal{D}(G)$  be the smallest value that the largest label can have.

For example, for the complete graph  $K_n$ , the labels must form a Sidon set. Hence,  $\mathcal{D}(K_n) = (1 + o(1))n^2$ . Rather surprisingly, we demonstrate that there are graphs with only  $n^{\frac{3}{2}+o(1)}$  edges achieving this bound.

More generally, we study the maximum value of  $\mathcal{D}(G)$  that a graph  $G$  of the given order  $n$  and size  $m$  can have. We obtain bounds which are sharp up to a logarithmic multiplicative factor. The analogous problem for pairwise sums is considered as well. Our results, in particular, disprove a conjecture of Wood.

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## 1. Introduction

Let  $G$  be a graph. A *difference-magic labelling* of  $G$  is an injective mapping  $l : V(G) \rightarrow \mathbb{N}$  (into positive integers) such that the  $e(G)$  numbers

$$|l(x) - l(y)|, \quad \{x, y\} \in E(G),$$

are pairwise distinct.

It is trivial to see that every graph admits a difference-magic labelling, so a natural question to ask is how economical it can be. More precisely, we should like to determine

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the *difference-magic number*  $\mathcal{D}(G)$  which is the smallest  $k$  such that a difference-magic labelling of  $G$  into  $[k] := \{1, \dots, k\}$  exists.

For example, it is easy to see that if  $G$  is the complete graph of order  $n$ , then  $\mathcal{D}(G)$  is precisely  $s_n$ , the smallest  $s$  such that  $[s]$  contains a *Sidon* subset of size  $n$ . (A set  $A \subset \mathbb{Z}$  is *Sidon* if all sums  $a + b$  with  $a, b \in A$  and  $a \leq b$  are distinct.) The latter problem is well studied; the results of Singer [13] and Erdős and Turán [8] (see e.g. Halberstam and Roth [10, Chapter II]) imply that  $s_n = (1 + o(1))n^2$ . Erdős [5] offered \$500 for proving or disproving that  $s_n = n^2 + O(n)$ .

Here we deal with

$$\mathcal{D}(n, m) := \max\{\mathcal{D}(G) : v(G) = n, e(G) = m\},$$

the maximum value of  $\mathcal{D}(G)$  for a graph  $G$  of order  $n$  and size  $m$ .

It turns out that

$$\mathcal{D}(n, m) = (1 + o(1))n^2 \quad \text{if } m / \sqrt{n^3 \ln n} \rightarrow \infty. \quad (1)$$

In fact, a random graph of order  $n$  with the appropriate edge probability demonstrates (1). We find it surprising that graphs so sparse (with only  $n^{\frac{3}{2} + o(1)}$  edges) have the  $\mathcal{D}$ -function asymptotically the same as that of the complete graph.

What happens for smaller  $m$ ? The obvious choice is to consider random graphs of suitable density. This, indeed, leads to interesting results. Let  $G \in \mathcal{G}(n, p)$ , that is,  $G$  is a random graph on  $n$  vertices where each edge is included in  $G$  independently of others and with probability  $p$ . If  $p = O((\ln n/n)^{1/2})$  and  $p > n^{-1+\varepsilon}$ , then

$$\mathcal{D}(G) = \Theta(n^3 p^2 / \ln n). \quad (2)$$

A lower bound on  $\mathcal{D}(n, m)$  can be obtained by adding isolated vertices to a random graph and figuring out the best parameters to choose. On the other hand, the simple labelling procedure described in Section 4 gives an upper bound that is within an  $O((\ln n)^{2/3})$ -factor of the lower bound. Roughly, we obtain

$$\mathcal{D}(n, m) = m^{4/3 + O\left(\frac{\ln \ln m}{\ln m}\right)}, \quad \text{if } m = O\left(\sqrt{n^3 \ln n}\right), \quad (3)$$

unless  $m = o(n^{3/4})$  when  $\mathcal{D}(n, m) = (1 + o(1))n$ . All details (with more precise expressions for the error terms) can be found in the corresponding sections.

Let us define a *sum-magic labelling* of a graph  $G$  as an injection  $l : V(G) \rightarrow \mathbb{N}$  such that all  $e(G)$  sums  $l(x) + l(y)$ ,  $\{x, y\} \in E(G)$ , are pairwise distinct. We ask for the *sum-magic number*  $\mathcal{S}(G)$ , the smallest value that the largest label can have, and for

$$\mathcal{S}(n, m) := \max\{\mathcal{S}(G) : v(G) = n, e(G) = m\}.$$

It is not surprising that most of the methods on the  $\mathcal{D}$ -function transfer to  $\mathcal{S}$ , giving similar bounds. (In particular, (3) holds for  $\mathcal{S}(n, m)$  as well.) However, there is one peculiar distinction. While Corollary 2 states that  $\mathcal{S}(K_n) = (1 + o(1))n^2$ , Theorem 3 shows that there is a constant  $c > 0$  such that  $\mathcal{S}(n, m) < (1 - c)n^2$  whenever  $m \leq cn^2$ . Random graphs are far worse in hitting  $(1 + o(1))n^2$ : this happens only when the random graph is almost complete.

Wood [15] defines an *edge-magic injection* with the *magic sum*  $s$  as an injection  $l : V(G) \cup E(G) \rightarrow \mathbb{N}$  such that for any edge  $\{a, b\} \in E(G)$  we have  $s = l(a) + l(b) + l(\{a, b\})$ . Let  $\mathcal{E}(G)$  be the smallest possible value of  $s$ . Wood [15, Section 7] conjectured that there is an absolute constant  $C$  such that for any graph  $G$  we have  $\mathcal{E}(G) \leq C(v(G) + e(G))$ . Clearly, the vertex labels of any edge-magic injection form a sum-magic labelling, so  $\mathcal{E}(G) \geq \mathcal{S}(G)$  and random graphs disprove Wood's conjecture.

One can also ask what is the value of, for example,

$$\mathcal{S}_{\min}(n, m) := \min\{\mathcal{S}(G) : v(G) = n, e(G) = m\}.$$

This is the inverse problem to maximising the number of distinct pairwise sums that a set  $A \subset [s]$  of given size  $n$  can have. This question is investigated by Pikhurko [12].

## 2. Some preliminary results

Let  $A \in \binom{[m]}{n}$ , meaning that  $A$  is an  $n$ -subset of  $[m]$ .

Recall that  $A$  is called a *Sidon set* if the sums  $a + b$ ,  $a, b \in A$  with  $a \geq b$ , are pairwise distinct, which is equivalent to all differences  $a - b$ ,  $a, b \in A$  with  $a > b$ , being pairwise distinct. Erdős and Turán [8] proved that this property implies that  $m \geq (1 + o(1))n^2$ . The following results show that, in a sense, it is the condition on differences (rather than that on sums) which pushes  $\max A$  upwards.

For  $i \in [m - 1]$  let  $g_i$  be the number of representations  $i = a - b$  with  $a, b \in A$ . Thus, if  $n^2 \geq (1 + \varepsilon)m$  (and  $n$  is large), then there must be  $i$  with  $g_i \geq 2$ . Although the following theorem strengthens this claim considerably, its proof goes via an easy modification of the original argument of Erdős and Turán [8]. A similar result (in a more precise form) was independently obtained by Ferrara, Kohayakawa and Rödl [9, Lemma 12].

Let  $f_+ = f$  if  $f > 0$  and  $f_+ = 0$  otherwise.

**Theorem 1.** *Let  $\varepsilon > 0$  be fixed and  $n \rightarrow \infty$ . Let  $A \in \binom{[m]}{n}$ . If  $n^2 \geq (1 + \varepsilon)m$ , then  $g = \Omega(n^2)$ , where  $g := \sum_{i=1}^{m-1} (g_i - 1)_+$ .*

**Proof.** Let  $t := cn^2$ , where  $c = c(\varepsilon) > 0$  is a small constant. Assume  $t \in \mathbb{N}$ . Define

$$A_i := A \cap [i, i + t - 1] \quad \text{and} \quad a_i := |A_i|, \quad i \in [2 - t, m],$$

where  $[i, j] := \{i, i + 1, \dots, j\}$ .

Let  $\mathcal{X}$  consist of all quadruples  $(a, b, i, x)$  such that  $x = a - b > 0$  and  $a, b \in A_i$ . Using the identity  $\sum_{i=2-t}^m a_i = nt$  and the quadratic-arithmetic mean inequality, we obtain

$$|\mathcal{X}| = \sum_{i=2-t}^m \binom{a_i}{2} = \frac{1}{2} \sum_{i=2-t}^m a_i^2 - \frac{nt}{2} \geq \frac{(nt)^2}{2(m+t-1)} - \frac{nt}{2}. \quad (4)$$

Each  $x \in [t - 1]$  is included in  $g_x \cdot (t - x) \leq (t - x) + t(g_x - 1)_+$  quadruples. Hence,

$$|\mathcal{X}| \leq \sum_{x=1}^{t-1} (t - x + t(g_x - 1)_+) = \frac{t(t-1)}{2} + gt. \quad (5)$$

By choosing  $c$  sufficiently small, we can ensure that the right-hand side of (4) is, for example, at least  $(1 + \frac{\varepsilon}{2})\frac{t^2}{2}$ , which together with (5) implies the theorem.  $\square$

We will need [Theorem 1](#) in [Section 3](#). Here we demonstrate another application.

**Corollary 2.**  $S(K_n) = (1 + o(1))n^2$ .

**Proof.** Let  $A$  be the label set of a sum-magic labelling. Note that  $A$  need not be Sidon as it may well happen that  $a - c = c - b$  for  $a, b, c \in A$ . However, if  $a - b = c - d$  with  $a \notin \{b, c\}$ , then either  $a = d$  or  $b = c$ . It follows that  $g_x \leq 2$  for any  $x > 0$  and, if  $g_x = 2$ , then there are  $a, b, c \in A$  with  $a - b = c - a = x$ . If  $a - b' = c' - a \neq 0$ , then we have  $b' + c' = 2a = b + c$  and thus  $\{b', c'\} = \{b, c\}$ . Hence, no  $a$  can appear for more than one  $x$  in the above manner. We conclude that  $g \leq |A|$ , implying the claim by [Theorem 1](#).  $\square$

The natural analogue of [Theorem 1](#) in terms of the number of solutions to  $x = a + b$ ,  $a, b \in A$ , is not true, as the following construction of Erdős and Freud [6] demonstrates. Let  $S \in \binom{[t]}{s}$  be a Sidon set with  $t = (1 + o(1))s^2$ . (Such sets were constructed by Singer [13].) Let  $X = S \cup S'$ , where

$$S' := 3t + 1 - S := \{3t + 1 - a : a \in S\} \subset [2t + 1, 3t].$$

Clearly,  $S + S \subset [2, 2t]$ ,  $S + S' \subset [2t + 2, 4t]$  and  $S' + S' \subset [4t + 2, 6t]$  are disjoint. Hence, all sums  $a + b$ ,  $a, b \in X$  with  $a \leq b$ , are pairwise distinct except those  $s$  sums which are equal to  $3t + 1$ . If the complement of an order- $n$  graph  $G$  has a matching covering all but  $r = o(n)$  vertices, then considering the first  $n$  elements of the set  $X$  constructed above for  $s := \frac{n+r}{2}$ , we conclude

$$S(G) \leq (3/4 + o(1))n^2. \quad (6)$$

By modifying the above construction, we can show one of the results claimed in the Introduction.

**Theorem 3.** *There is a constant  $c > 0$  such that if  $m \leq cn^2$ , then*

$$S(n, m) \leq (1 - c)n^2. \quad (7)$$

**Proof.** Let  $\alpha = 0.9$ , for example. In the above construction of  $X = S \cup S'$  let  $Y \subset X$  consist of the first  $n := \lfloor (1 + \alpha)s \rfloor$  elements of  $X$ . As it was shown by Erdős and Freud [6, Lemma 1], any asymptotically maximum Sidon subset of  $[t]$  is almost uniformly distributed. This implies that  $\max Y = (2 + \alpha + o(1))t$ .

Now, all sums in  $Y + Y$  are distinct except those sums which equal  $3t + 1$ . The number of these exceptional sums is  $\lfloor \alpha s \rfloor = (\frac{\alpha}{1+\alpha} + o(1))n$ . So, if the complement of an order- $n$  graph  $G$  has a matching of size bigger than  $0.48n > (\frac{\alpha}{1+\alpha} + o(1))n$ , then

$$S(G) \leq (2 + \alpha + o(1))t = \frac{2 + \alpha + o(1)}{(1 + \alpha)^2} n^2 < 0.9n^2.$$

It follows from the Tutte 1-factor theorem [14] that a matching of size  $0.48n$  in the complement  $\overline{G}$  is guaranteed if  $e(G) \leq \delta n^2$  for some constant  $\delta > 0$ . Now, the theorem follows.  $\square$

**Remark.** Random graphs do not provide good examples if we want to achieve  $\mathcal{S}(G) = (1+o(1))n^2$ : this happens only when  $1-p = O(\frac{\ln n}{n})$ . Indeed, Erdős and Rényi [7] (cf. Bollobás and Thomason [3]) showed that if  $p \leq 1 - (1+\varepsilon)\frac{\ln n}{2n}$  then with high probability the complement of  $G \in \mathcal{G}(n, p)$  has an almost perfect matching; so then (6) holds.

### 3. Random graphs

**Theorem 4.** Fix any  $\delta > 0$ . Let  $G \in \mathcal{G}(n, p)$ , where  $n \rightarrow \infty$  and  $p \in (0, 1)$  is a function of  $n$  such that  $np/\ln n \rightarrow \infty$ . Let  $\lambda := p\sqrt{n/\ln n}$ . Then almost surely  $\mathcal{D}(G) \geq d$  and  $\mathcal{S}(G) \geq s$ , where

$$d := \begin{cases} (1-\delta)n^2, & \text{if } \lambda \rightarrow \infty, \\ \left(\frac{\lambda^2}{16+2\lambda^2} - \delta\right)n^2, & \text{if } \lambda = \Theta(1), \\ \left(\frac{1}{16} - \delta\right)\frac{n^3 p^2}{\ln n}, & \text{if } \lambda = o(1). \end{cases} \quad s := \begin{cases} \left(\frac{1}{4} + \frac{1}{(\pi+2)^2} - \delta\right)n^2, & \text{if } \lambda \rightarrow \infty, \\ \left(\frac{\lambda^2}{32+4\lambda^2} - \delta\right)n^2, & \text{if } \lambda = \Theta(1), \\ \left(\frac{1}{32} - \delta\right)\frac{n^3 p^2}{\ln n}, & \text{if } \lambda = o(1). \end{cases}$$

**Proof.** We prove the lower bound on  $\mathcal{D}(G)$ . Let  $[n]$  be the vertex set. Let  $\varepsilon > 0$  be a small constant depending on  $\delta$ . Assume  $d \in \mathbb{N}$ .

Fix an injective mapping  $l : [n] \rightarrow [d]$ . Now, let us choose  $G \in \mathcal{G}(n, p)$ . We want to bound the probability  $p'$  that all differences  $l(i) - l(j)$ , with  $\{i, j\} \in \binom{[n]}{2}$  being an edge of  $\mathcal{G}(n, p)$  and  $l(i) > l(j)$ , are pairwise distinct. If  $u$  is an upper bound on  $p'$  for any  $l$ , then the probability that  $G \in \mathcal{G}(n, p)$  satisfies  $\mathcal{D}(G) \leq d$  is at most  $n! \binom{d}{n} u < d^n u$ . Hence, if we can show that  $p' = o(d^{-n})$ , then almost surely  $\mathcal{D}(G) > d$ .

For  $k \in [d]$ , let  $g_k$  be the number of representations  $k = l(i) - l(j)$  with  $i, j \in [n]$ . Let  $t := \binom{n}{2} = \sum_{k=1}^d g_k$ . Clearly,

$$p' = \prod_{k=1}^d p_k, \quad (8)$$

where  $p_k = (1-p)^{g_k} + g_k p(1-p)^{g_k-1}$  is the probability of selecting at most one edge with difference  $k$ . (Note that the formula is also valid for  $g_k = 0$  and  $g_k = 1$ , when  $p_k = 1$ .) It is routine to see that

$$p' = \prod_{k=1}^d p_k \leq ((1-p)^{t/d} + (pt/d)(1-p)^{(t/d)-1})^d. \quad (9)$$

**Case 1.**  $p = o(\sqrt{\ln n/n})$ , that is,  $\lambda = o(1)$ .

We have  $t/d \rightarrow \infty$  and  $pt/d = o(1)$ . Using the inequality  $e^{-x}(1+x) \leq 1 - (\frac{1}{2} - \varepsilon)x^2$  valid if  $x > 0$  is small, we deduce from (9) the required bound on  $p'$ :

$$\begin{aligned} p' &\leq ((1-p)^{t/d}(1+pt/d+2p^2t/d))^d \leq (e^{-pt/d}(1+pt/d+2p^2t/d))^d \\ &\leq (1 - (1/2 - \varepsilon)(pt/d)^2 + 2p^2t/d)^d \leq e^{-(1/2-2\varepsilon)(pt)^2/d} = o(e^{-n \ln d}). \end{aligned}$$

**Case 2.**  $p = \Theta(\sqrt{\ln n/n})$ , that is,  $\lambda = \Theta(1)$ .

We have  $t/d = O(1)$  so we can simply take the Taylor expansion of (9) to obtain the required bound:

$$p' \leq \left(1 + \left(\frac{t}{2d} - \frac{t^2}{2d^2}\right)p^2 + O(p^3)\right)^d \leq e^{\frac{1}{2}(t - \frac{t^2}{d} + \varepsilon)p^2} = o(e^{-n \ln d}).$$

**Case 3.**  $p\sqrt{n/\ln n} \rightarrow \infty$ , that is,  $\lambda \rightarrow \infty$ .

By [Theorem 1](#) we know that  $g := \sum_{k=1}^d (g_k - 1)_+ = \Omega(n^2)$ . It is routine to see that if  $g_i \geq g_j + 2$ , then the right-hand side of (8) increases if we replace  $g_i$  and  $g_j$  by  $g_i - 1$  and  $g_j + 1$  respectively. Hence,

$$p' \leq ((1 - p)^2 + 2p(1 - p))^g = (1 - p^2)^g = o(d^{-n}),$$

as required.

Let us turn to the sum-magic number. Fix an injection  $l : [n] \rightarrow [s]$ . For  $k \in [2s]$  define  $g_k$  as the number of representations  $k = l(i) + l(j)$  with  $1 \leq i < j \leq n$ . Let  $t := \binom{n}{2} = \sum_{k=1}^{2s} g_k$ . The remainder of the proof goes via the obvious modification of the argument for  $\mathcal{D}(G)$  except that for  $\lambda \rightarrow \infty$  we use the result by Pikhurko [[12](#), Theorem 2] which implies that  $\sum_{k=1}^{2s} (g_k - 1)_+ = \Omega(n^2)$ . (If we are content with  $s = (\frac{1}{4} - \delta)n^2$ , then  $\sum_{k=1}^{2s} (g_k - 1)_+ = \Omega(n^2)$  follows by trivial counting.) The reader should have little difficulty in filling in all missing details.  $\square$

**Remark.** There is a jump in the lower bounds when we change from the case  $\lambda = \Theta(1)$  to  $\lambda \rightarrow \infty$ . It should be possible to ‘smoothen’ this by improving our bounds for large but bounded  $\lambda$ . However, the calculations seem to be rather unpleasant, so we do not go into the details.

**Remark.** As it was mentioned in the introduction, [Theorem 4](#) disproves the conjecture of Wood in view of the inequality  $\mathcal{E}(G) \geq \mathcal{S}(G)$ . Indeed, if we take  $G \in \mathcal{G}(n, n^{-1/2})$  for example, then almost surely  $e(G) = (\frac{1}{2} + o(1))n^{3/2}$  while  $\mathcal{D}(G) = \Omega(n^2/\ln n)$ . With a bit of extra work it is possible to show that under the assumptions of [Theorem 4](#) we have almost surely  $\mathcal{E}(G) \geq 2s$ . To do this, prove that, almost surely, any sum-magic labelling of  $G$  has  $\Omega(n)$  labels which are greater than  $s$  and there is an edge connecting two such labels. We leave the details to the interested reader.

Now let us turn to upper bounds.

**Theorem 5.** Let  $\delta > 0$  be fixed. Let  $G \in \mathcal{G}(n, p)$ , where  $n \rightarrow \infty$  and  $p \in (0, 1)$  is a function of  $n$  such that  $\frac{\ln(np)}{\ln n} \rightarrow \infty$  and  $p = O((\ln n/n)^{1/2})$ . Then almost surely  $\mathcal{D}(G) \leq 2m$  and  $\mathcal{S}(G) \leq m$ , where

$$m := (1 + \delta) \frac{n^3 p^2}{\ln(np)}. \quad (10)$$

**Proof.** Let us estimate  $\mathcal{S}(G)$ . (The case of  $\mathcal{D}(G)$  is dealt with almost identically.)

We can assume that  $\delta$  is sufficiently small and  $m \in \mathbb{N}$ . Let  $n$  be large and  $\varepsilon > 0$  be a small constant depending on  $\delta$ . Let  $V(G) = [n]$  be the vertex set. Chernoff’s bound [[4](#)]

implies that almost surely we have

$$|\Gamma(i+1) \cap [i]| - ip \leq \varepsilon np, \quad \text{for all } i \in [0, n-1], \quad (11)$$

where  $\Gamma(i+1)$  is the set of neighbours of  $i+1 \in V(G)$ .

Consider the conditional distribution of  $G$  given (11). We have gained the very useful control over the edges while some important properties of  $G \in \mathcal{G}(n, p)$  are preserved. (That is, almost sure events stay so; the random set  $\Gamma(i+1) \cap [i]$  is independent from  $G[[i]]$ , etc.)

We choose vertex labels one by one, doing the label arithmetic in  $M = \mathbb{Z}/m\mathbb{Z}$  (that is, modulo  $m$ ). Our labelling  $l : V(G) \rightarrow [m]$  will have the property that the sums  $l(x) + l(y)$ ,  $\{x, y\} \in E(G)$ , will be pairwise distinct modulo  $m$ .

Suppose that we have already chosen labels for the vertices in  $I := [i]$ .

Let

$$K := \{l(x) + l(y) : \{x, y\} \in E(G), x, y \in I\} \subset M,$$

and  $k := |K|$ . By (11),

$$k \leq \left(\frac{1}{2} + \varepsilon\right) inp. \quad (12)$$

Clearly, we can find a suitable label for  $i+1$  if

$$M \setminus l(I) \not\subset \bigcup_{x \in I \cap \Gamma(i+1)} (K - l(x)), \quad (13)$$

that is, if the translates  $K - l(x)$ ,  $x \in I \cap \Gamma(i+1)$ , do not cover  $M \setminus l(I)$ .

This is obviously the case if

$$|M \setminus (\cup_{x \in I} (K - l(x)))| \geq n,$$

so let us assume otherwise. Then we have  $m - ik \leq n$ , which implies by (12) that

$$i \geq n\sqrt{2p/\ln(np)}. \quad (14)$$

Now, we have to overcome the difficulty that  $i$  is large enough to potentially refute (13). In outline, we fix the labelling  $l$  of  $I$  and then choose the random set  $I \cap \Gamma(i+1)$ . The labels  $l(x)$ ,  $x \in I \cap \Gamma(i+1)$ , are random variables. If the translates  $K - l(x)$  cover the whole of  $M \setminus l(I)$ , then for every  $z \in M \setminus l(I)$  at least one element  $l(x) \in K - z$  is chosen. We prove that this is unlikely.

Let  $S$  consist of those elements from  $M \setminus l(I)$  which are covered by at most

$$t := \lfloor (1 + \varepsilon)kn/m \rfloor$$

of the translates  $K - l(x)$ ,  $x \in I$ . Clearly,

$$|M \setminus (S \cup l(I))| \times (1 + \varepsilon)kn/m \leq kn,$$

so

$$s := |S| \geq \frac{\varepsilon m}{1 + \varepsilon} - n \geq \varepsilon m/2.$$

Let  $\gamma' = \lfloor ip + \varepsilon np \rfloor$  and  $\gamma = \lceil ip + 2\varepsilon np \rceil$ . Let us choose  $y_1, \dots, y_\gamma \in I$ , one by one, independently and uniformly distributed. Of course, some of these might coincide. Let  $Y := \{y_1, \dots, y_\gamma\}$ , ignoring multiple occurrences of the same vertex. The probability that  $|Y| \leq \gamma'$  is at most

$$\begin{aligned} \binom{i}{\gamma'} \left( \frac{\gamma'}{i} \right)^\gamma &\leq \left( \frac{ei}{ip + \varepsilon np} \right)^{ip + \varepsilon np} \left( \frac{ip + \varepsilon np}{i} \right)^{ip + 2\varepsilon np} \\ &= e^{ip + \varepsilon np} \left( p + \frac{\varepsilon np}{i} \right)^{\varepsilon np}. \end{aligned}$$

The last expression, as a function of a real-valued argument  $i \geq 0$ , is first decreasing and then increasing in  $i$  so it is maximised if either  $i = n$  or  $i$  achieves the lower bound (14). In either case, the result can be bounded by  $o(n^{-1})$ . Thus, the set  $Y$  has at least  $\gamma'$  elements with probability  $1 - o(n^{-1})$ .

Let the random variable  $U$  count the number of  $x \in S$  which belong to none of  $K - l(y)$ ,  $y \in Y$ .

We consider the martingale  $(U_0, \dots, U_\gamma)$ , where  $U_j$  is the expected value of  $U$  after having exposed the first  $j$  vertices  $y_1, \dots, y_j$ . Clearly, each new vertex changes  $U$  by at most  $k$ .

It is easy to estimate  $U_0$ , the expectation of  $U$ :

$$E(U) \geq |S| \frac{\binom{i-t}{\gamma}}{\binom{i}{\gamma}} \geq s \left( \frac{i-t-\gamma+1}{i-\gamma+1} \right)^\gamma \geq s e^{-(1+\varepsilon)\gamma t/i} \geq s(np)^{-\frac{1}{2} + \frac{\delta}{4}}.$$

(Note that  $t = o(i)$  by the definition of  $t$  and  $\gamma = o(i)$  by (14).)

By applying the Hoeffding–Azuma inequality [2, 11] (see e.g. Alon and Spencer [1, Theorem 7.2.1]) we obtain

$$\begin{aligned} \Pr\{U = 0\} &\leq \Pr\{|U - E(U)| \geq E(U)\} \leq \exp\left(-\frac{(E(U))^2}{2k^2\gamma}\right) \\ &= \exp(-\Omega((np)^{\frac{\delta}{2}}/\ln^2(np))) = o(n^{-1}). \end{aligned}$$

Hence, the event that  $|Y| < \gamma'$  for some  $i$  or  $U = 0$  has probability  $o(1)$ . Of course, when we select a random  $a$ -subset of  $Y$ , we obtain a uniformly distributed  $a$ -subset of  $I$ . Note that  $|\Gamma(i+1) \cap [i]| \leq \gamma'$  by (11). We can find a distribution for  $a \in [0, i]$  such that when we first choose  $a$ , then  $Y$  as above, then a random  $a$ -subset of  $Y$ , we obtain precisely the distribution of  $\Gamma(i+1) \cap [i]$ , conditioned on (11).

Hence, almost surely for any  $i$ , condition (13) holds; that is, we can always choose an appropriate label.  $\square$

#### 4. General graphs

Let us prove upper bounds that apply to arbitrary graphs. The obvious greedy algorithm gives the following (cf. Wood [15, Theorem 4]).



**Lemma 6.** For any graph  $G$  we have

$$\mathcal{D}(G) \leq 2\Delta(G)e(G) + v(G),$$

$$\mathcal{S}(G) \leq \Delta(G)e(G) + v(G).$$

**Proof.** Let us bound  $\mathcal{S}(G)$ , for example. We choose vertex labels one by one. When we consider a vertex  $i \in V(G)$ , we are forbidden to choose a previously used label as well as any number of the form  $l(u) + l(v) - l(w)$  where  $\{u, v\}, \{w, i\} \in E(G)$  and the labels of  $u, v$  and  $w$  have already been chosen. This forbids at most  $v(G) - 1 + d(i)e(G)$  elements so we can always proceed.  $\square$

Remarkably, the trivial Lemma 6 is not far from the truth: if applied to  $G \in \mathcal{G}(n, p)$ , with  $\frac{np}{\ln n} \rightarrow \infty$ , it gives a bound within the multiplicative factor of  $O(\ln n)$  from the actual value. It is an interesting open problem to determine the maximum value of  $\mathcal{D}(G)$  and  $\mathcal{S}(G)$  over all graphs of order  $n$  and maximum degree at most  $d$ .

For the functions  $\mathcal{D}(n, m)$  and  $\mathcal{S}(n, m)$  we obtain the following upper bounds.

**Theorem 7.** Let  $n \rightarrow \infty$  and  $m \leq \binom{n}{2}$ . Then

$$\mathcal{D}(n, m) \leq n + (2^{4/3} + o(1))m^{4/3},$$

$$\mathcal{S}(n, m) \leq n + (2^{2/3} + o(1))m^{4/3}.$$

**Proof.** Let us deal with  $\mathcal{D}(n, m)$  here. Let  $n$  be large and  $G$  be an arbitrary graph of order  $n$  and size  $m$ . It is easy to see that  $\mathcal{D}(n, m) = n$  if  $m = O(1)$ , so assume that  $m \rightarrow \infty$ .

Order the vertices of  $G$  by their degrees:  $d(x_1) \geq \dots \geq d(x_n)$ . Let  $k = \lceil (2m)^{2/3} \rceil$ . Label vertices  $x_1, \dots, x_k$  by a Sidon  $k$ -subset of  $[s]$ ,  $s = \lfloor (1 + \varepsilon)k^2 \rfloor$ . We try to label the remaining vertices one by one using labels from  $[n + s]$ . When choosing a label for  $x_i$ , the forbidden values are the already assigned labels,  $i - 1$  of them, as well as the numbers  $l(x_u) \pm (l(x_v) - l(x_w))$ , where  $u, v, w \in [i - 1]$  with  $\{x_i, x_u\}, \{x_v, x_w\} \in E(G)$ , at most  $2md(x_i)$  numbers. But  $d(x_j) \geq d(x_i)$  for any  $j < i$ , hence  $d(x_i) \leq \frac{2m}{i} \leq \frac{2m}{k}$  and the total number of forbidden labels is at most

$$n - 1 + \frac{4m^2}{k} < n + s;$$

that is, we can always find a suitable label.  $\square$

Needless to say, we have a trivial upper bound, namely  $(1 + o(1))n^2$ .

Good lower bounds on  $\mathcal{D}(n, m)$  and  $\mathcal{S}(n, m)$  are provided by random graphs plus isolated vertices. Our aim is to choose  $v \leq n$  such that, if we define  $p$  by  $p \binom{v}{2} = (1 - \varepsilon)m$ , the bound of Theorem 4 for  $G \in \mathcal{G}(v, p)$  is as large as possible. In order not to clutter this paper with details we compute only the order of magnitude, not bothering about multiplicative constants.

If  $m = \Omega(n^{3/2}\sqrt{\ln n})$ , then we take  $v = n$ . Almost surely  $\mathcal{D}(G), \mathcal{S}(G) = \Omega(n^2)$ . Otherwise, take  $v = \Theta(m^{2/3}(\ln m)^{-1/3}) < n$ . Now, the lower bound is

$$\mathcal{D}(n, m), \mathcal{S}(n, m) = \Omega(m^{4/3}(\ln m)^{-2/3}), \quad \text{for } m = o(n^{3/2}\sqrt{\ln n}).$$

Also, note the trivial lower bound  $\mathcal{D}(n, m), \mathcal{S}(n, m) \geq n$ .

A little more careful analysis shows that there is an absolute constant  $C$  such that our lower and upper bounds on  $\mathcal{D}(n, m)$  and  $\mathcal{S}(n, m)$  are within factor  $C(\ln n)^{2/3}$  for any  $m, n$ . This poses an intriguing problem of closing this gap.

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